

How Well Do Traffic Engineering Objective Functions Meet TE Requirements ?

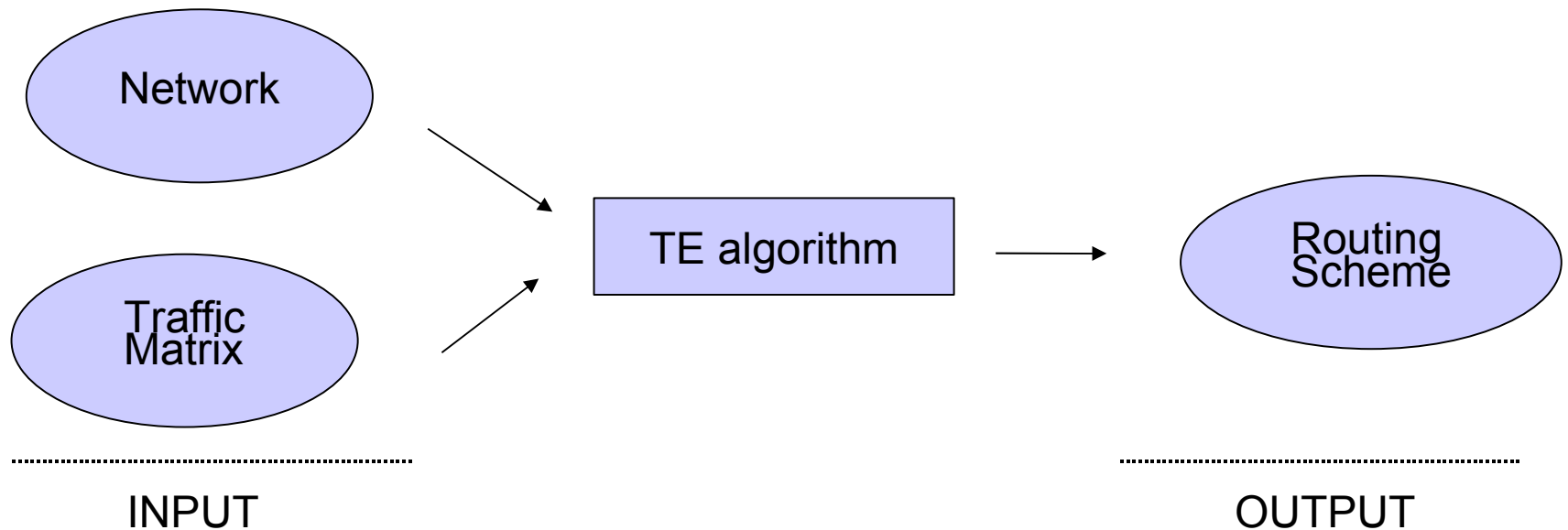
Research Unit in Networking (RUN)
Université de Liège (ULg, Belgium)

- **Presentation** : S. Balon
- **Authors** : S. Balon, F. Skivée, G. Leduc

- Introduction
- TE requirements & Notations
- TE Metrics
- TE objective functions
- Simulation Results
- Conclusion

- **Introduction**
- TE requirements & Notations
- TE Metrics
- TE objective functions
- Simulation Results
- Conclusion

- Goal of the study: Compare and evaluate how TE objective functions fulfil TE requirements
- We consider the "routing problem" :



- TE algorithm is supposed to find a "good" routing scheme (i.e. a set of paths)

- What is a "good" routing scheme?

It is generally a routing scheme that associates a high score to a chosen *objective function*

- We can find many objective functions which are very different in the literature. How is it possible to choose the best ?
- We also want to separate the benefits due to heuristic from the benefit due to the objective function of well known TE algorithms presented as a whole solution in many papers

- But how well do objective functions reflect TE requirements?
- We try to answer this question with simulations. We find the optimum for each Objective Function with LP formulation of the routing problem. We arrive to this kind of simulation results :

Objective Function	u_{max}	θ_{tot}	...
<i>MIRA</i>	100%	6504 Mbps	...
<i>Delay</i>	42%	6309 Mbps	...
⋮	⋮	⋮	⋮

- Introduction
- **TE requirements & Notations**
- TE Metrics
- TE objective functions
- Simulation Results
- Conclusion

- TE objectives (view of IETF):
 - performance optimization of operational networks
 - minimize packet loss, delay, delay variation
 - maximize throughput
 - enforce Service Level Agreements
 - minimize maximum resource utilisation
- How do well-known TE algorithms address these TE objectives?

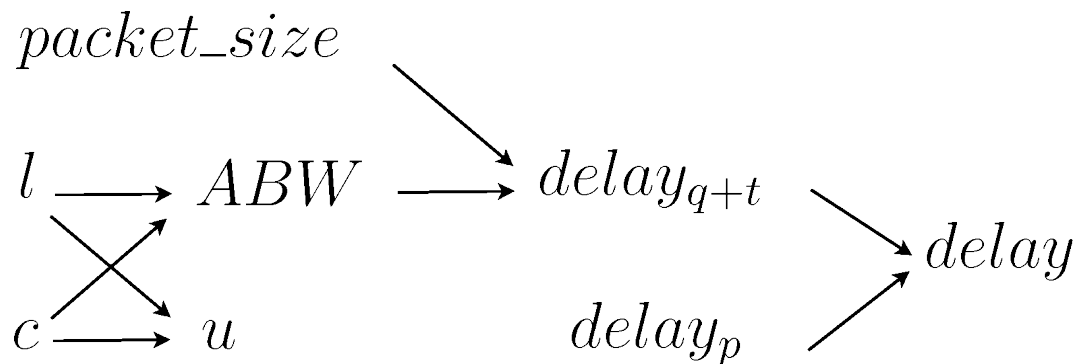
- Notations :
 - Network modelled by a graph $G=(N,A)$
 - Each link a has a capacity $c(a)$
 - Traffic Matrix D
 - Each link a has a load $l(a)$
 - Utilisation of link a is $u(a) = l(a) / c(a)$
 - Available bandwidth of link a is $ABW(a) = c(a) - l(a)$

- Introduction
- TE requirements & Notations
- **TE Metrics**
- TE objective functions
- Simulation Results
- Conclusion

- Mean queuing + transmission delay :

$$Delay_{q+t} = \frac{mean_packet_size}{ABW}$$

- Link parameters:



- Residual Max Flow between nodes s and t : θ_{st}
- TE Metric Summary:

	Metric characterising good current state	Metric characterising likely good future
Link _(a)	$Delay_a$	u_a, ABW_a
Path _(s,t)	$\sum_{a \in \mathcal{P}(s,t)} Delay_a$	$\theta_{st}, \max_{a \in \mathcal{P}(s,t)} u_a$
Network	$\frac{\sum_{a \in A} Delay_a}{ A },$ $\frac{\sum_{a \in A} l_a \cdot Delay_a}{AllTr}$	$\min_{(s,t)} \theta_{st}, \max_{a \in A} u_a$ $\sum_{(s,t)} \theta_{st}$

- Introduction
- TE requirements & Notations
- TE Metrics
- **TE objective functions**
- Simulation Results
- Conclusion

- **InvCap:** $\sum_{a \in A} u_a$ *This function is minimized by CISCO recommended IP routing – which computes shortest paths based on link cost = $1/c(a)$*
- **U_{\max} :** u_{\max} *This function minimizes the maximal link utilisation.*
- **MinHop:** $\sum_{a \in A} l_a$ *This function is minimised by minimum hop routing (if the used paths are the path with minimum number of links)*

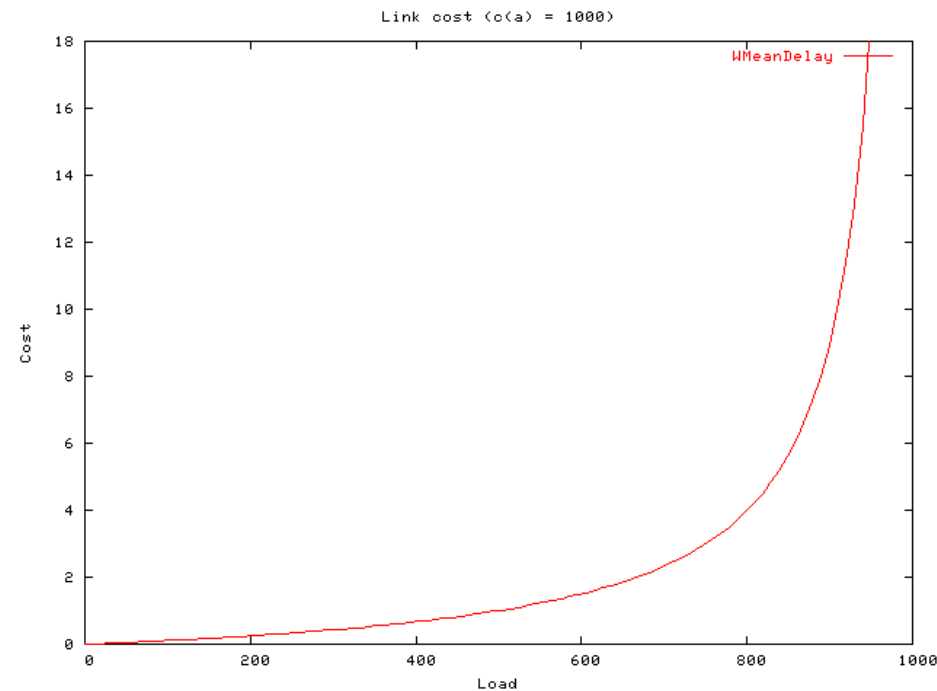
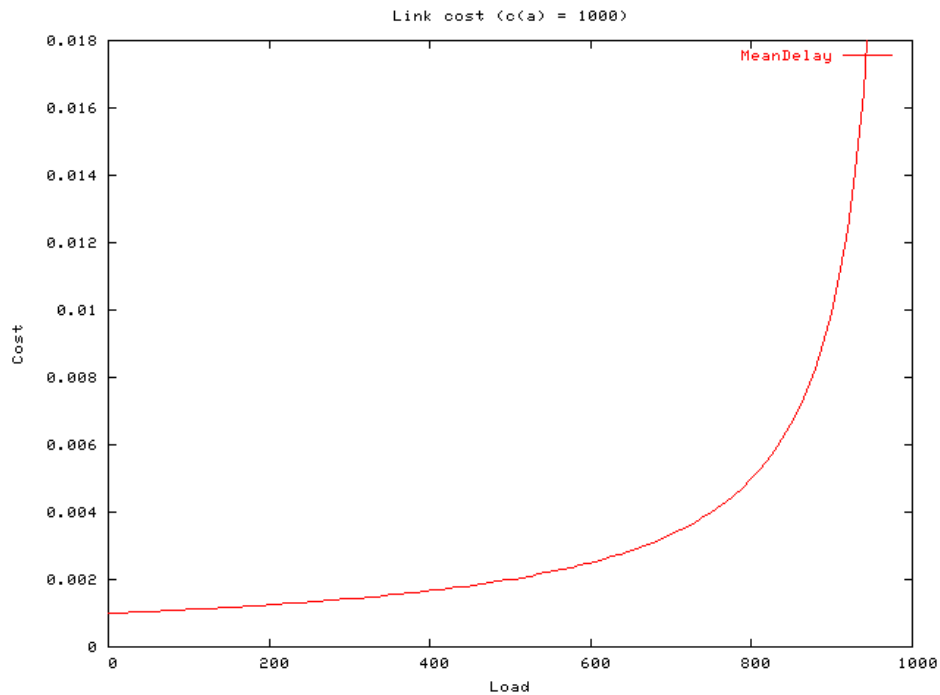
Transmission + Queuing delay of link $a \approx \frac{1}{c(a) - l(a)}$

- MeanDelay: $\sum_{a \in A} \frac{1}{c(a) - l(a)}$ *Minimizes the mean link delay*

- WMeanDelay: $\sum_{a \in A} \frac{l(a)}{c(a) - l(a)} = \sum_{a \in A} \frac{1}{\frac{1}{u(a)} - 1}$

Minimizes the weighted mean link delay = weighted mean path delay

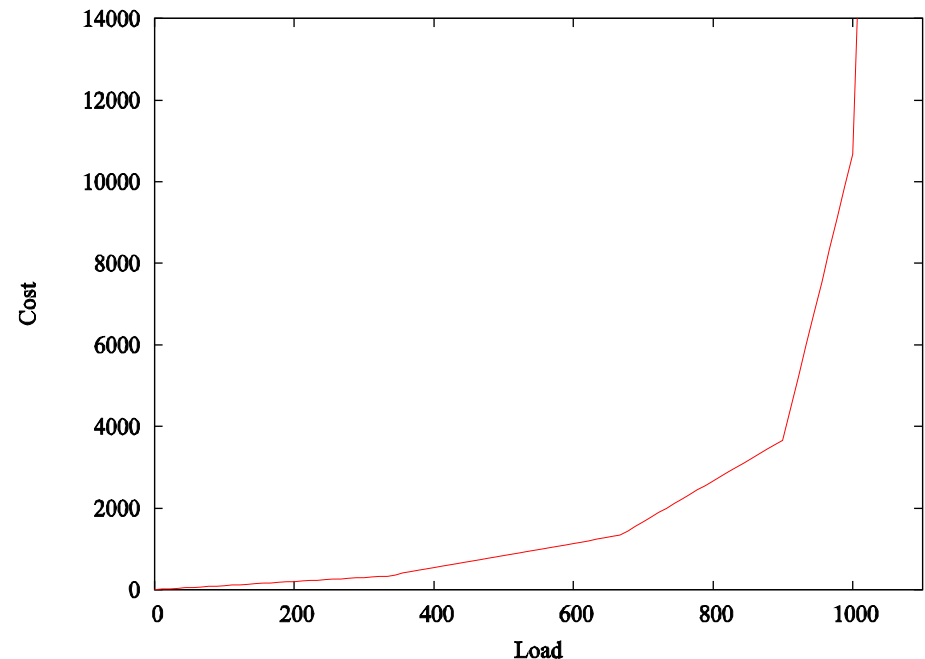
- Delay link cost:



- Fortz: $\min \phi = \sum_{a \in A} \phi_a$

ϕ_a

(for a link of capacity 1000)



*Associates a high cost to high loaded links
so that paths should avoid these links*

- MIRA: concept of minimum interference routing
- If θ_{st} is the max flow from s to t on the residual topology, $MIRA = \max \sum_{(s,t)} \alpha_{st} \theta_{st}$
- α_{st} are administrative weights

Try to minimize the decrease in the total residual max flow due to the routing of some trafic on a path

- Degrande: $C_B.B - C_U.U$
- Balance: $B = 1 - u_{max}$
- Network utilisation: $U = \sum_{a \in A} u(a)$

Try to minimize both the maximal link utilisation and the mean link utilisation

- *Blanchy* :

$$\underbrace{\sum_{a \in A} (u(a) - u_{mean})^2}_{\text{Load balancing term}} + \alpha \underbrace{\sum_{a \in A} (u(a))^2}_{\text{Traffic minimisation term}}$$

The load balancing term try to flatten the link utilisations while the traffic minimization term try to limit the size of paths which would otherwise make to long detours.

- Introduction
- TE requirements & Notations
- TE Metrics
- TE objective functions
- **Simulation Results**
- Conclusion

- Simulation to compare all presented objective functions:
 - LP formulation of the routing problem
 - Linearize non-linear objective functions
 - Solve the problem to optimality
 - Compare all the results
- All the objectives functions are compared with the same algorithm → easy to compare
- Results on three different networks (→ one random network with random Traffic Matrix and two real networks with real Traffic Matrices)

- Results on a network of 25 nodes (random Waxman topology):

Objective function	u_{max} %		u_{per10} %		u_{mean} %		ABW_{min} Mbps		l_{mean} Mbps		θ_{tot} Mbps	
	TM	2TM	TM	2TM	TM	2TM	TM	2TM	TM	2TM	TM	2TM
<i>Fortz</i>	(1.14)	(1.14)	(1.28)	(1.33)	(1.26)	(1.21)	(0.67)	(0.55)	(1.03)	(1.05)	(0.97)	(0.95)
<i>MIRA</i>	<i>100</i>	<i>100</i>	(1.40)	(1.48)	(1.17)	(1.16)	<i>0.0</i>	<i>0.0</i>	(1.15)	(1.10)	6504	5012
<i>Blanchy</i>	(1.22)	(1.23)	26.0	50.0	(1.13)	(1.12)	(0.88)	531	(1.12)	(1.11)	(0.96)	(0.94)
<i>Delay</i>	(1.20)	(1.08)	(1.17)	(1.20)	(1.04)	(1.11)	882.0	(0.95)	(1.16)	(1.11)	(0.97)	(0.95)
<i>InvCap</i>	<i>(2.07)</i>	<i>100</i>	(1.55)	(1.61)	15.7	31.5	882.0	<i>0.0</i>	(1.21)	(1.20)	(0.98)	(0.96)
u_{max}	34.9	69.7	(1.15)	(1.20)	(1.07)	(1.12)	(0.74)	(0.57)	(1.17)	(1.11)	(0.97)	(0.95)
<i>Degrande</i>	34.9	69.7	(1.35)	(1.39)	(1.05)	(1.05)	(0.74)	(0.57)	(1.19)	(1.18)	(0.97)	(0.95)
<i>MinHop</i>	<i>100</i>	<i>100</i>	(1.29)	(1.43)	(1.27)	(1.25)	<i>0.0</i>	<i>0.0</i>	781	1578	(0.97)	(0.95)

The table contains absolute optimal values (in green, bold, without parentheses), or relative non-optimal values (between parentheses) with respect to the optimal one. The values that are less than 10% from the optimal value are bold. Finally values that are 2 times worse than the optimum are in italic and red. For each metric we show the values for the actual traffic matrix (TM) and for the doubled traffic matrix (2TM).

- Metrics at Low Load (LL) and High Load (HL)

Objective Function	u_{max}		u_{per10}		u_{mean}		ABW_{min}		l_{mean}		θ_{tot}	
	LL	HL	LL	HL	LL	HL	LL	HL	LL	HL	LL	HL
<i>Fortz</i>	✓	✓	✓	✓	✓	✓	±	±	✓	✓	✓	✓
<i>MIRA</i>	•	•	•	•	✓	✓	•	•	✓	✓	✓✓	✓✓
<i>Blanchy</i>	✓	•	✓	✓	✓	✓	✓	•	✓	✓	✓	✓
<i>Delay</i>	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓
<i>InvCap</i>	•	•	✓	±	✓✓	✓✓	±	•	✓	✓	✓	✓
<i>Degrande</i>	✓✓	✓✓	✓	✓	✓	✓	✓	✓	✓	✓	✓	✓

- Introduction
- TE requirements & Notations
- TE Metrics
- TE objective functions
- Simulation Results
- **Conclusion**

- We have shown how well-known network-wide TE objective functions reflect TE requirements
- Our results show that they are not equivalent
- Although delay_{q+t} is negligible, this Objective Function gives good results for all TE metrics
- Best results for *Delay*, *Degrade* and *Fortz* (in that order)
- Objective basis to choose the best TE function

- Contact : *Simon.Balon@ulg.ac.be*

- TOTEM project

<http://totem.info.ucl.ac.be/>

<http://totem.run.montefiore.ulg.ac.be/>



- E-NEXT NoE

<http://www.ist-e-next.net/>



- FNRS - Belgian National Fund for Scientific Research

<http://www.fnrs.be/>



FNRS